

Disjoint sets

CS 491 – Competitive Programming

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Objectives

- ▶ Use the up-tree data structure to model disjoint sets
- ▶ Use path compression to make this structure efficient

What do these groups have in common?

- ▶ Lutherans
- ▶ Jacobites
- ▶ Marxists
- ▶

Disjoint Sets

Properties of disjoint sets

- ▶ If a is in the same set as b , then everything in the same set as a is in the same set as b .
- ▶ Want to declare that a and b are in the same set.
- ▶ Want to check if a and b are in the same set.
- ▶ Want to ask how many elements are in the set.
- ▶ Want to check if all known elements are in the set.

And we want to do all of these things very very quickly.

- ▶ The datastructure for this is the up-tree.

The idea

Representation

- ▶ Keep a vector ds representing all your elements.
- ▶ A positive entry means this element is in the same element as the referent.
- ▶ A negative entry means this element is the root. The magnitude is the number of elements.

Initialization

- ▶ To start, initialize the vector with all -1s.

Find operation

- ▶ $\text{find}(i) = i$, if $ds[i] < 0$
- ▶ $\text{find}(i) = \text{find}(ds[i])$, otherwise.

Add

Relating a and b .

- ▶ Let $sa = \text{find}(a)$
- ▶ Let $sb = \text{find}(b)$

Choices:

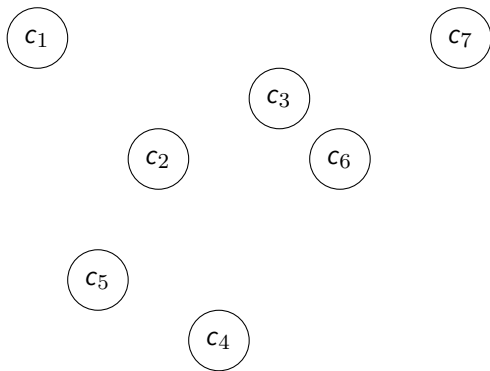
- ▶ $sa = sb$, already in the same set.
- ▶ $ds[sa] > ds[sb]$, then sb is the larger set.
 - ▶ Let $ds[sb] = ds[sb] + ds[sa]$ and $ds[sa] = sb$
- ▶ Symmetrically for the other case.

Code

```
int find(int c, vi ds) {
    if (ds[c] < 0)
        return c;
    else
        return find(ds[c], ds);
}
```

```
int add(int c1, int c2, vi ds) {
    int sa = find(c1);
    int sb = find(c2);
    if (sa == sb)
        return sa;
    if (ds[sa] > ds[sb]) {
        ds[sb] += ds[sa]; ds[sa] = sb; return sb;
    } else {
        ds[sa] += ds[sb]; ds[sb] = sa; return sa;
    }
}
```

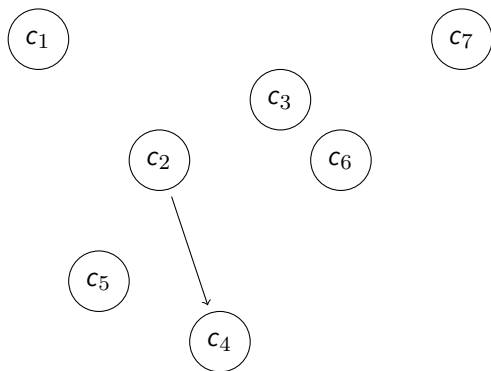
Example, pt 1



1	2	3	4	5	6	7
-1	-1	-1	-1	-1	-1	-1

▶ next connect 2 and 4

Example, pt 2

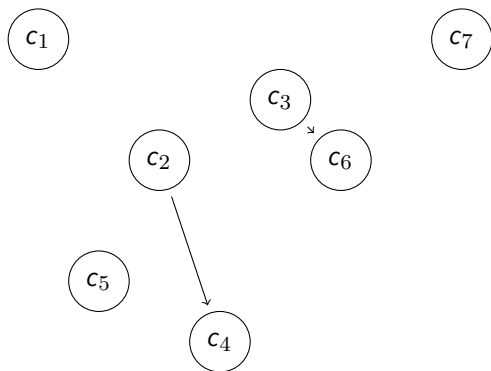


- ▶ connect 2 and 4

1	2	3	4	5	6	7
-1	4	-1	-2	-1	-1	-1

- ▶ next connect 3 and 6

Example, pt 3

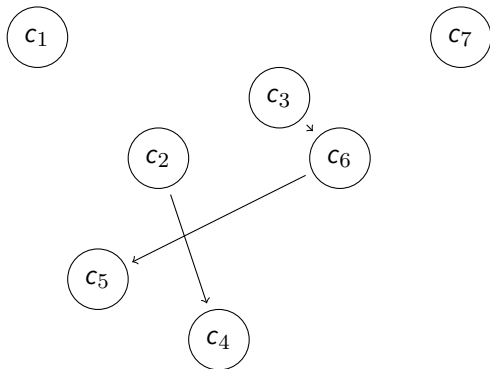


- ▶ connect 3 and 6

1	2	3	4	5	6	7
-1	4	6	-2	-1	-2	-1

- ▶ next connect 6 and 5

Example, pt 4

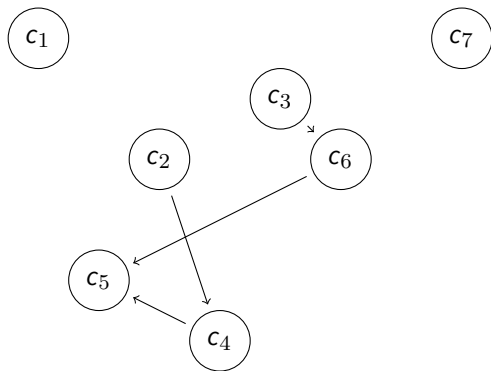


- ▶ connect 6 and 5

1	2	3	4	5	6	7
-1	4	6	-2	-3	5	-1

- ▶ next connect 2 and 3

Example, pt 5

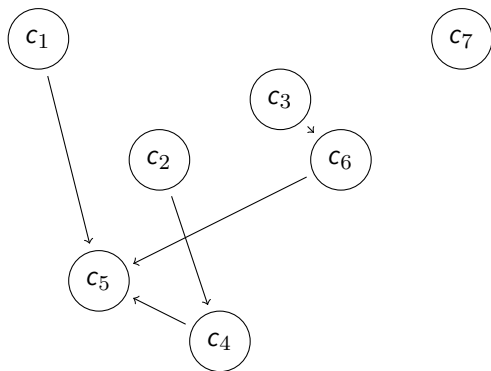


- ▶ connect 2 and 3 (will connect 4 to 5)

1	2	3	4	5	6	7
-1	4	6	5	-5	5	-1

- ▶ next connect 3 and 1

Example, pt 6

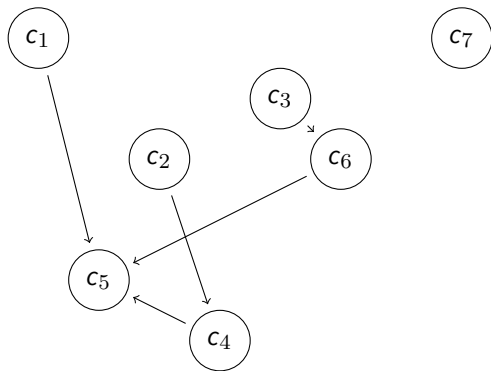


► connect 3 and 1

1	2	3	4	5	6	7
5	4	6	5	-6	5	-1

Path Compression

- ▶ One problem is that these paths can get quite long if we are unlucky.
- ▶ Solution: update `find` so it compresses the paths.

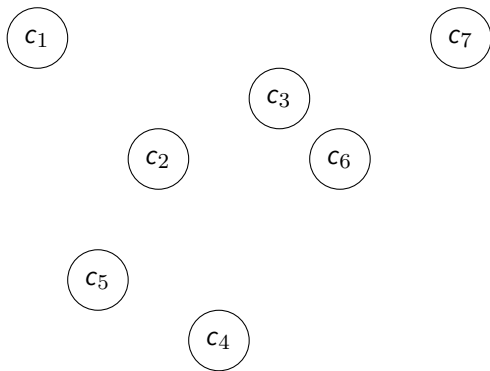


Find with Path Compression

```
int find(int c, vi ds) {
    if (ds[c] < 0)
        return c;
    else
        return ds[c] = find(ds[c], ds);
}
```

```
int add(int c1, int c2, vi ds) {
    int sa = find(c1);
    int sb = find(c2);
    if (sa == sb)
        return sa;
    if (ds[sa] > ds[sb]) {
        ds[sb] += ds[sa]; ds[sa] = sb; return sb;
    } else {
        ds[sa] += ds[sb]; ds[sb] = sa; return sa;
    }
}
```

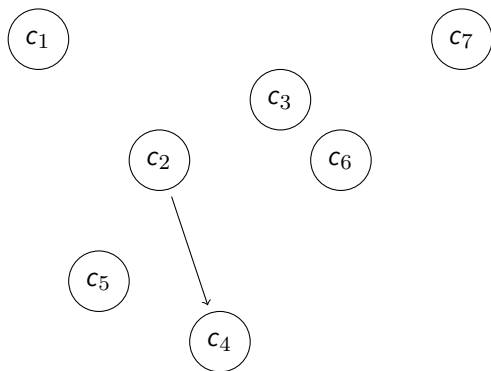
Example, pt 1



1	2	3	4	5	6	7
-1	-1	-1	-1	-1	-1	-1

▶ next connect 2 and 4

Example, pt 2

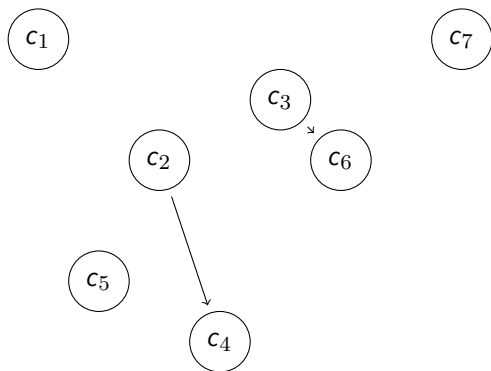


- ▶ connect 2 and 4

1	2	3	4	5	6	7
-1	4	-1	-2	-1	-1	-1

- ▶ next connect 3 and 6

Example, pt 3

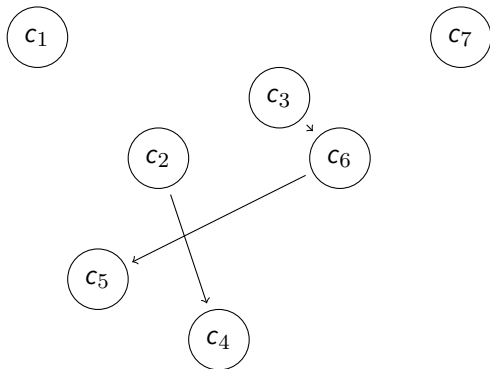


- ▶ connect 3 and 6

1	2	3	4	5	6	7
-1	4	6	-2	-1	-2	-1

- ▶ next connect 6 and 5

Example, pt 4

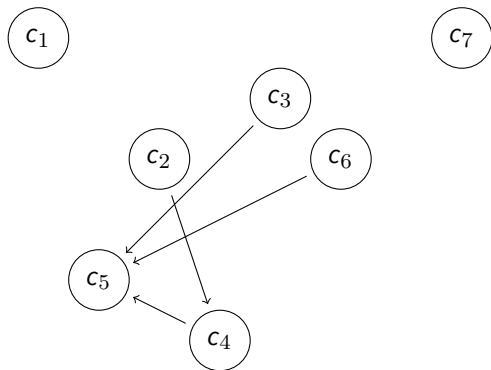


- ▶ connect 6 and 5

1	2	3	4	5	6	7
-1	4	6	-2	-3	5	-1

- ▶ next connect 2 and 3

Example, pt 5

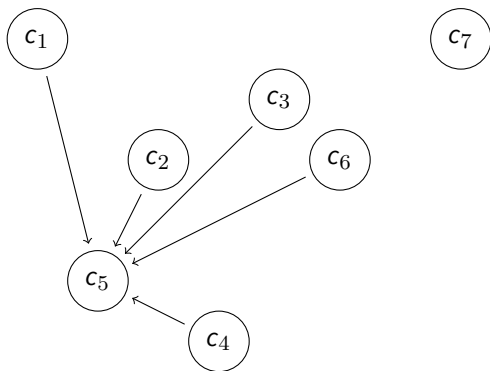


- ▶ connect 2 and 3 (will connect 4 to 5, and update 3)

1	2	3	4	5	6	7
-1	4	5	5	-5	5	-1

- ▶ next connect 3 and 1

Example, pt 6



► connect 3 and 1

1	2	3	4	5	6	7
5	5	5	5	-6	5	-1