

# Segment Trees

CS 491 – Competitive Programming

**Dr. Mattox Beckman**

UNIVERSITY OF ILLINOIS AT URBANA-CHAMPAIGN  
DEPARTMENT OF COMPUTER SCIENCE

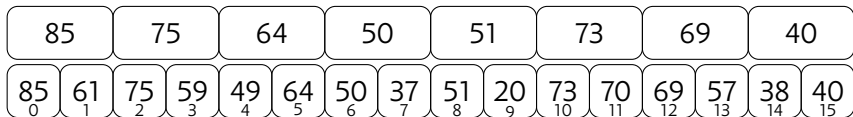
Fall 2023

# Running Example

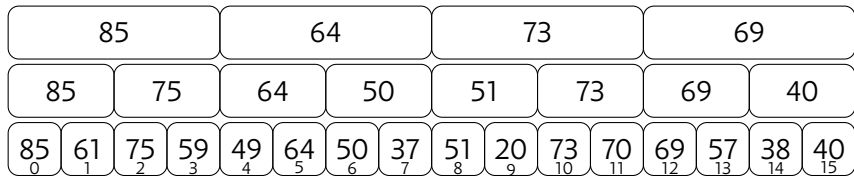
▶ Consider the following array:

|    |    |    |    |    |    |    |    |    |    |    |    |    |    |    |    |
|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|----|
| 85 | 61 | 75 | 59 | 49 | 64 | 50 | 37 | 51 | 20 | 73 | 70 | 69 | 57 | 38 | 40 |
| 0  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | 9  | 10 | 11 | 12 | 13 | 14 | 15 |

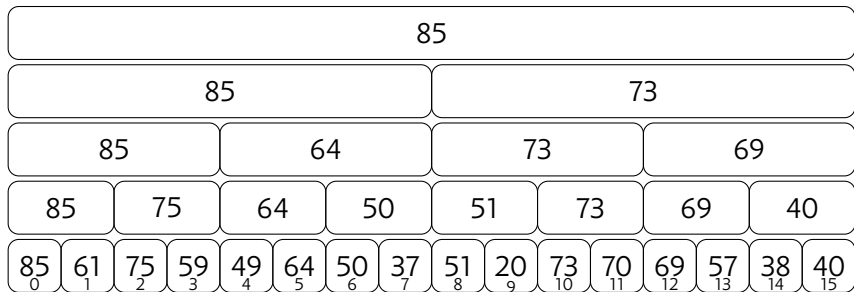
# Segment Trees, Level 1



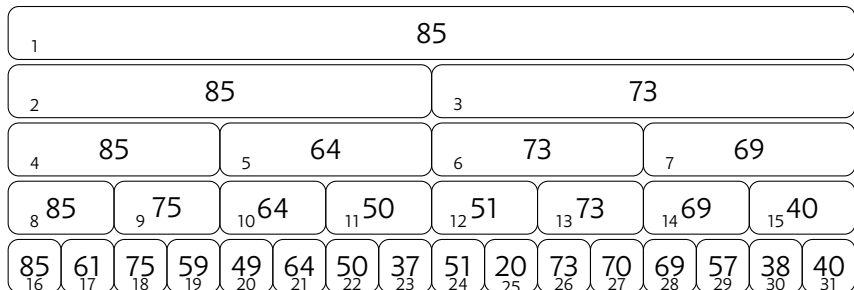
## Segment Trees, Level 2



## Segment Trees, Level 3 and 4



# Segment Trees, Numbering the Elements



## Build the Segement Tree

- ▶ L and R give you the bounds with respect to the original array.
- ▶ This code gives you a min range.

```
1 void build(int p, int L, int R) {
2     if (L == R) // as L == R, either one is fine
3         st[p] = data[L]; // store the data
4     else {
5         // recursively compute the values
6         build(left(p) , L , (L + R) / 2);
7         build(right(p), (L + R) / 2 + 1, R );
8         int p1 = left(p), p2 = right(p);
9         st[p] = min(st[p1],st[p2]);
10    }
```

## Query the Tree

- ▶ L and R give you the bounds with respect to the original array.
- ▶ i and j give you the bounds for the query

```
1  int rmq(int p, int L, int R, int i, int j) {
2      if (i > R || j < L) return -1; // current segment outside
3      if (L >= i && R <= j) return st[p];
4      // compute the min position in the left and right part of
5      int lm = rmq(left(p), L, (L+R)/2, i, j);
6      int rm = rmq(right(p), (L+R)/2+1, R, i, j);
7      if (lm == -1) return rm;
8      // if we try to access segment outside query
9      if (rm == -1) return lm;
10     return min(lm,rm);
11 }
```

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