

Sqrt Decomposition

CS 491 – Competitive Programming

Dr. Mattox Beckman

UNIVERSITY OF ILLINOIS AT URBANA-CHAMPAIGN
DEPARTMENT OF COMPUTER SCIENCE

Fall 2024

Objectives

- ▶ Use sqrt decomposition to improve the time complexity of large problems.

Running Example

- ▶ Consider the following array:

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
85	61	75	59	49	64	50	37	51	20	73	70	69	57	38	40

- ▶ What is an algorithm, given i and j , of returning the sum between these numbers (inclusive)?
- ▶ What is an algorithm, given i and j , of returning the max between these numbers (inclusive)?

Code for Sum

```
1  vi run;
2  int a = 0;
3
4  run.push_back(0); // sentinel
5  for(int i: data) {
6      a += data;
7      run.push_back(a);
8  }
9
10 int sum(int i, int j) {
11     return run[j+1] - run[i];
12 }
```

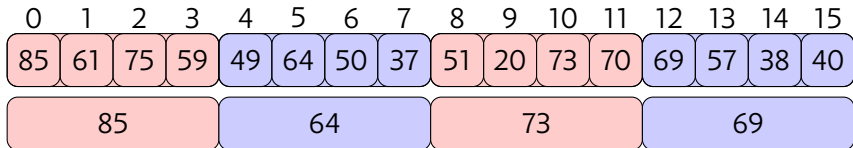
- ▶ We can't do a "running max" though.

Kotlin Version

```
1 val size = readln().toInt()
2 val data = readln().split(' ').map { it.toInt() }
3 val run = data.runningFold(0) { acc, num ->
4     acc + num }.drop(1)
5
6 fun sum(run : List<Int>, i : Int, j : Int) : Int {
7     return run[j+1] - run[i]
8 }
```

Solution

- ▶ We can create a separate array to handle each block of \sqrt{n} size.



- ▶ What is the max number between...
 - ▶ 0 and 2?
 - ▶ 4 and 11?
 - ▶ 3 and 8?
- ▶ What is the resulting time complexity?

Preprocessing Code

- ▶ `sq` contains the sqrt decompositions.
- ▶ `data` contains the raw data.

```
1 vi data, sq;
2 int n,a,d;
3
4 cin >> n;
5 int s = sqrt(n);
6 sq = vi(s+1);
7
8 for(int i=0,j=0; i < s && j < n; ++i ) {
9     cin >> d; data.push_back(d);
10    sq[i] = d;
11    for(k=1; k<s; ++j, ++k) {
12        cin >> d; data.push_back(d);
13        sq[i] = max(sq[i],d);
14    }
15 }
```

Preprocessing Code, Kotlin

```
1 import kotlin.math.sqrt
2
3 val size = readln().toInt()
4 val data : List<Int> = readln().split(' ').map {
5     it.toInt() }
6
7 val nsq = sqrt(data.size.toDouble()).toInt()
8 val chunk = data.chunked(nsq)
9 val sq = chunk.map { sublist -> sublist.max() }
```


Search

Note this ignores some edge cases

```
1  int findMax(int i, int j, vi &data, vi &sq, int s) {
2      int a;
3      a = data[i];
4      while (i % s > 0 && i <= j) {
5          a = max(a,data[i]);
6          ++i;
7      }
8
9      while ( (j+1) % s > 0) {
10         a = max(a,data[j]);
11         --j;
12     }
13
14     for(k = i/s; k<=j/s; ++k)
15         a = max(a,sq[k]);
16 }
```

Kotlin Version

```
1 fun findMax(i: Int , j : Int, data : List<Int>, sq : List<Int>) {
2     var a = data[i]
3     val iup = minOf(j, i + s - (i % s) - 1)
4     val jdown = j - (j % s)
5     val b1 = i / s
6     val b2 = j / s
7
8     if (b1 == b2) { // Same block
9         return maxOf(a,data.slice(i..j).max())}
10    if (i % s > 0) {a = maxOf(a,data.slice(i..iup).max())}
11    if (jdown <= j) {a = maxOf(a,data.slice(jdown..j).max())}
12    if (b1+1 < b2) {a = maxOf(a,sq.slice(b1+1..b2-1).max())}
13    return a
14 }
```

15

16