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Traveling Sales Person CS 491 – Competitive Programming

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Objectives

- Understand how to use DP to solve the TSP problem.
- Use the bitmask technique to represent the cities already visited.
- Know the limit of *n* for brute force and DP problems.

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The Problem

- You are given a set of nodes with weighted edges
 - cities and cost / time for travel
- Want to make a tour, visit all cities, return to start.
- What is the cheapest way to do this?
- Time complexity
 - $\mathcal{O}(n!)$ check each permutation
 - Fix first city to take advantage of symmetry gives (n 1)! solutions.
 - In a contest, brute force check can work up to about n = 11, max.

Setup

```
int main() {
   cin >> n;
   vvi adj(n);
   for(i=0; i<n; ++i)</pre>
      for(j=0; j<n; ++j) {</pre>
          cin >> c;
          adj[i].push back(c);
      }
   mx = (1 << n) - 1;
   cout << "Best path has cost " << tsp(adj,mx,0,1) << end</pre>
}
```

Using DP

```
vvi dp(16,vi(65536)); int n;
int tsp(vvi &costs, int &mx, int cur, int state) {
   if (dp[cur][state]>0) return dp[cur][state];
   if (state == mx) return costs[cur][0]; // return home
   int minleft = INF; int bit=2;
  for(int i=1; i<n; ++i) {</pre>
     if ( (state & bit) == 0) { // i not visited
       minleft = min(minleft, costs[cur][i] +
                          tsp(costs,mx,i,state | bit));
    }
    bit <<= 1:
  }
  return dp[cur][state]=minleft;
}
```